

Lecture 4: System of Equations

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1 Direct Methods

1.1 Gaussian Elimination

1.2 Gauss Jordan Algorithm

2 Matrix Factorisation

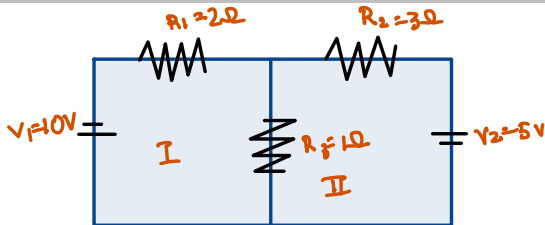
2.1 LU Decomposition

2.2 LDL & Cholesky Decomposition



Direct Methods

Direct Methods Matrix Factorisation



Kirchoff's Voltage law: $V_i = I_i R_i$ (Ohm's law) for $i=1,2,3$

In loop 1

$$10 - 2I_1 - (I_1 - I_2) = 0$$

$$\Rightarrow -3I_1 + I_2 = -10$$

In loop 2,

$$5 - 3I_2 - (I_2 - I_1) = 0$$

$$\Rightarrow I_1 - 4I_2 = -5$$

} I_1, I_2



Gaussian Elimination

Direct Methods Matrix Factorisation

$$\left. \begin{array}{l} R_1 \\ R_2 \\ \vdots \\ R_n \end{array} \right\} \begin{array}{l} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\ \vdots \\ a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n \end{array} \quad \left. \begin{array}{l} \\ \\ \\ \end{array} \right\} \{x_i\}_{i=1}^n \rightarrow \text{Unknowns}$$

Row numbers.

1. Scalar multiplication : $R_i \rightarrow \lambda R_i$, for λ real number $\neq 0$
2. Row replacement : $R_i \rightarrow R_i + \lambda R_j$ for $j \neq i$
3. Row interchange : $R_i \leftrightarrow R_j$, $i \neq j$

1. Augmented Matrix

$$\underline{A}\underline{x} = \underline{b}; \quad A = \{a_{ij}\}; \quad \underline{b} = \{b_i\}_{i=1}^n, \quad \underline{x} = \{x_i\}_{i=1}^n$$

$$[A, \underline{b}] = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} \left| \begin{array}{c} b_1 \\ b_2 \\ \vdots \\ b_n \end{array} \right. = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & a_{1,n+1} \\ a_{21} & a_{22} & \dots & a_{2n} & a_{2,n+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} & a_{n,n+1} \end{bmatrix}$$

$a_{1,n+1}$ (pointing to b_1) and $a_{n,n+1}$ (pointing to b_n)



Gaussian Elimination

Direct Methods Matrix Factorisation

$$\begin{array}{l} R_1 \\ R_2 \\ \vdots \\ R_n \end{array} \left[\begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & a_{1,n+1} \\ a_{21} & a_{22} & \dots & a_{2n} & a_{2,n+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} & a_{n,n+1} \end{array} \right]$$

How?



$$\left[\begin{array}{cccc|c} \bar{a}_{11} & \bar{a}_{12} & \dots & \bar{a}_{1n} & \bar{a}_{1,n+1} \\ 0 & \bar{a}_{22} & \dots & \bar{a}_{2n} & \bar{a}_{2,n+1} \\ 0 & 0 & \bar{a}_{33} & \bar{a}_{3n} & \bar{a}_{3,n+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \bar{a}_{nn} & \bar{a}_{n,n+1} \end{array} \right]$$

$$R_2 \rightarrow R_2 - \frac{a_{21}}{a_{11}} R_1 \quad (\text{so that } a_{21}=0)$$

$$\left[\begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & a_{1,n+1} \\ 0 & \bar{a}_{22} & \dots & \bar{a}_{2n} & \bar{a}_{2,n+1} \\ a_{31} & a_{32} & \dots & a_{3n} & a_{3,n+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} & a_{n,n+1} \end{array} \right]$$

$$R_j^0 \rightarrow R_j^0 - \frac{a_{j1}^0}{a_{11}} R_1, \quad j=2, \dots, n$$

my 1st col is

$$\left[\begin{array}{c} a_{11} \\ 0 \\ 0 \\ \vdots \\ 0 \end{array} \right]$$

Moving in this fashion, we get ←



Gaussian Elimination

Direct Methods Matrix Factorisation

$$\left[\begin{array}{cccc|c} \bar{a}_{11} & \bar{a}_{12} & \dots & \bar{a}_{1n} & \bar{a}_{1,n+1} \\ 0 & \bar{a}_{22} & \dots & \bar{a}_{2n} & \bar{a}_{2,n+1} \\ 0 & 0 & \bar{a}_{33} & \bar{a}_{3n} & \bar{a}_{3,n+1} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \bar{a}_{nn} & \bar{a}_{n,n+1} \end{array} \right] \sim$$

$$\begin{aligned} \bar{a}_{11}x_1 + \bar{a}_{12}x_2 + \dots + \bar{a}_{1n}x_n &= \bar{a}_{1,n+1} \\ \bar{a}_{22}x_2 + \dots + \bar{a}_{2n}x_n &= \bar{a}_{2,n+1} \\ &\vdots \\ \bar{a}_{n-1,n-1}x_{n-1} + \bar{a}_{n-1,n}x_n &= \bar{a}_{n-1,n+1} \\ \bar{a}_{nn}x_n &= \bar{a}_{n,n+1} \end{aligned}$$

$$x_n = \frac{\bar{a}_{n,n+1}}{\bar{a}_{nn}}$$

$$\bar{a}_{n-1,n-1}x_{n-1} + \bar{a}_{n-1,n}x_n = \bar{a}_{n-1,n+1}$$

$$x_{n-1} = \frac{\bar{a}_{n-1,n+1} - \bar{a}_{n-1,n}x_n}{\bar{a}_{n-1,n-1}}$$

We continue this to get $\{x_i\}_{i=1}^n$

Backward Substitution



Gaussian Elimination

Direct Methods Matrix Factorisation

Example: $[A, b] = \left[\begin{array}{ccc|c} 1 & 1 & 1 & 6 \\ 2 & 3 & 1 & 11 \\ 3 & 2 & 2 & 13 \end{array} \right]$ $R_2 \rightarrow R_2 - 2R_1$ $R_3 \rightarrow R_3 - 3R_1$

$$= \left[\begin{array}{ccc|c} 1 & 1 & 1 & 6 \\ 0 & 1 & -1 & -1 \\ 0 & -1 & -1 & -5 \end{array} \right]$$

$$R_3 \rightarrow R_3 + R_2 = \left[\begin{array}{ccc|c} 1 & 1 & 1 & 6 \\ 0 & 1 & -1 & -1 \\ 0 & 0 & -2 & -6 \end{array} \right]$$

With Backward subs: $\underline{x} = [1 \ 2 \ 3]^T$

Suppose is 0:

Example: $A = \left[\begin{array}{cccc|c} 1 & -1 & 2 & -1 & -8 \\ 2 & -2 & 3 & -3 & -20 \\ 1 & 1 & 1 & 0 & -2 \\ 1 & -1 & 4 & 3 & 4 \end{array} \right] \sim \left[\begin{array}{ccc|c|c} 1 & -1 & 2 & -1 & -8 \\ 0 & 0 & -1 & -1 & -4 \\ 0 & 2 & -1 & 1 & 6 \\ 0 & 0 & 2 & 4 & 12 \end{array} \right]$

Pivoting to avoid division by zero and then Row interchange; $R_2 \leftrightarrow R_3$

$$\left[\begin{array}{cc|cc} a_{11} & a_{12} & a_{13} & b_1 \\ 0 & 0 & a_{23} & b_2 \\ 0 & 0 & a_{33} & b_3 \end{array} \right] \rightarrow \left[\begin{array}{cc|cc} a_{11} & a_{12} & a_{13} & b_1 \\ 0 & 0 & a_{33} & b_3 \\ 0 & 0 & a_{23} & b_2 \end{array} \right]$$

$a_{33} = 0$ $b_3 \neq 0$

Locate the 1st non-zero entry

System no solution
or
 ∞ solution



Computational Complexity

Direct Methods Matrix Factorisation

Matrix-Reduction

$$A/S: \frac{n^3 - n}{3}$$

$$M/D: \frac{2n^3 + 3n^2 - 5n}{6}$$

$$\rightarrow \frac{n^3}{3} + n^2 - \frac{n}{3}$$

Backward Substitution

$$A/S: \frac{n^2 - n}{6}$$

$$M/D: \frac{n^2 + n}{2}$$

$$\rightarrow \frac{n^3}{3} + \frac{n^2}{2} - \frac{5n}{6}$$



Gauss Jordan Algorithm

Direct Methods Matrix Factorisation

$$\left[\begin{array}{cccc|c} \bar{a}_{11} & \bar{a}_{12} & \dots & \bar{a}_{1n} & \bar{a}_{1,n+1} \\ 0 & \bar{a}_{22} & \dots & \bar{a}_{2n} & \bar{a}_{2,n+1} \\ 0 & 0 & \bar{a}_{33} & \bar{a}_{3n} & \bar{a}_{3,n+1} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \bar{a}_{nn} & \bar{a}_{n,n+1} \end{array} \right] \rightarrow \left[\begin{array}{cccc|c} \hat{a}_{11} & 0 & \dots & 0 & \hat{a}_{1,n+1} \\ 0 & \hat{a}_{22} & \dots & 0 & \hat{a}_{2,n+1} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \dots & \hat{a}_{nn} & \hat{a}_{n,n+1} \end{array} \right]$$

$$x_i^0 = \frac{\hat{a}_{i,n+1}}{\hat{a}_{ii}}, \text{ for } i=1, \dots, n$$



LU Decomposition

Direct Methods Matrix Factorisation

L: Lower triangular matrix

U: Upper triangular matrix

$L\underline{x} = \underline{b}$ or $U\underline{x} = \underline{b} \rightarrow$ Complexity is $O(n^2)$

$$A = LU \rightarrow A\underline{x} = \underline{b}$$

$$\Rightarrow L(U\underline{x}) = \underline{b}$$

$$\Rightarrow L\underline{y} = \underline{b} \quad \text{where } \underbrace{U\underline{x} = \underline{y}}$$

Step 1: Solve this compute y

Step 2: Solve for x



LU Decomposition

Direct Methods Matrix Factorisation

Assumption: $a_{ii}^{(k)} \neq 0 \forall i$, and steps k . (No pivoting)

$$A \xrightarrow{\text{Gaussian elimination}} U \rightarrow \text{Compute } L \quad \begin{bmatrix} a_{11} \\ a_{21} \\ \vdots \\ a_{n1} \end{bmatrix} \times \begin{bmatrix} x \\ \vdots \\ x \end{bmatrix} \rightsquigarrow \begin{bmatrix} a_{11} \\ 0 \\ \vdots \\ 0 \end{bmatrix} \times \begin{bmatrix} x \\ \vdots \\ x \end{bmatrix}$$

$$R_j \rightarrow R_j - m_{j1} R_1, \quad \text{where } m_{j1} = \frac{a_{j1}^{(1)}}{a_{11}^{(1)}}, \quad j=2, \dots, n$$

equivalent to multiplying A with $M^{(1)}$ Same

$$M^{(1)} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ -m_{21} & 1 & 0 & \dots & 0 \\ -m_{31} & 0 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ -m_{n1} & 0 & 0 & \dots & 1 \end{bmatrix}$$

$$M^{(1)} A = \begin{bmatrix} 1 & 0 \\ -m_{21} & 1 \\ -m_{31} & 0 \\ \vdots & \vdots \\ -m_{n1} & 0 \end{bmatrix} \times \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \\ \vdots & \vdots \\ a_{n1} & a_{n2} \end{bmatrix} \times \begin{bmatrix} x \\ \vdots \\ x \end{bmatrix}$$

$$-m_{21} a_{11} + a_{21}$$

$$= -\frac{a_{21}}{a_{11}} a_{11} + a_{21} = 0$$



LU Decomposition

Direct Methods Matrix Factorisation

$$\underline{A}\underline{x} = \underline{b} \rightarrow \underbrace{M^{(1)}} A \underline{x} = M^{(1)} \underline{b} \rightarrow \text{First Gaussian elimination}$$

$$\begin{bmatrix} a_{11} & a_{12} & \dots \\ 0 & a_{22} & \dots \\ 0 & a_{32} & \dots \\ \vdots & \vdots & \ddots \\ 0 & a_{m2} & \dots \end{bmatrix}$$

$$M^{(2)} = \begin{bmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ 0 & -m_{32} & 1 & \dots \\ \vdots & \vdots & \vdots & \ddots \\ 0 & -m_{m2} & 0 & \dots \end{bmatrix}$$

$$\underbrace{M^{(2)} M^{(1)} A}_{A^{(3)}} \underline{x} = M^{(2)} M^{(1)} \underline{b}$$

Continuity

where $A^{(n)} = M^{(n-1)} M^{(n-2)} \dots M^{(1)} A = U$

What kind of matrix? Upper U \rightarrow Inverse L

$$L^{-1} A = U$$

$$A = LU$$



LU Decomposition

Direct Methods Matrix Factorisation



Definition (Diagonally Dominant Matrix)

A matrix $\mathbf{A} = (a_{ij})$ is said to be *diagonally dominant* if

$$|a_{ii}| \geq \sum_{\substack{j=1 \\ j \neq i}}^n |a_{ij}|, \quad i = 1, 2, \dots, n.$$

If the inequality is strict for every i , then \mathbf{A} is called *strictly diagonally dominant*.



Theorem

A strictly diagonally dominant matrix \mathbf{A} is nonsingular. Moreover, Gaussian elimination applied to a system $\mathbf{Ax} = \mathbf{b}$ can be carried out without row or column interchanges, and the computation is stable with respect to the growth of round-off errors.



Definition (Positive Definite Matrix)

A matrix \mathbf{A} is said to be *positive definite* if

$$\mathbf{x}^T \mathbf{A} \mathbf{x} > 0 \quad \text{for all } \mathbf{x} \neq \mathbf{0}.$$

If \mathbf{A} is symmetric and positive definite, it is called a *symmetric positive definite (SPD)* matrix.

